



Introduction to Parsing Part II



In Class

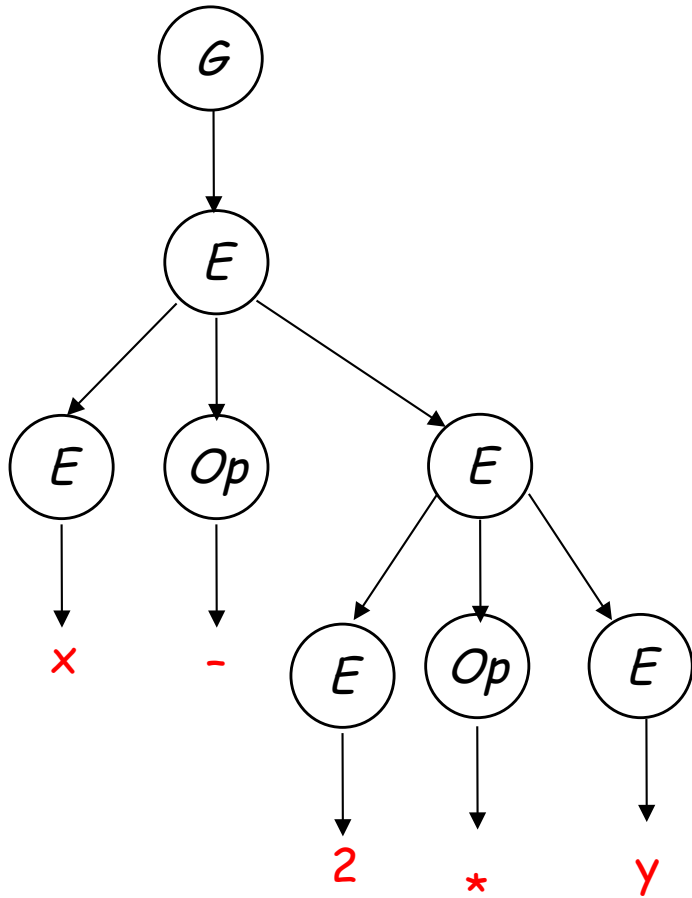
1	<i>Expr</i>	→	<i>Expr Op Expr</i>
2			<u>number</u>
3			<u>id</u>
4	<i>Op</i>	→	+
5			-
6			*
7			/

Produce a table showing the rightmost derivation for the equation below. Include in the first column the rule used and the second column the sentential form.

$$\underline{x} - \underline{2} * \underline{y}$$

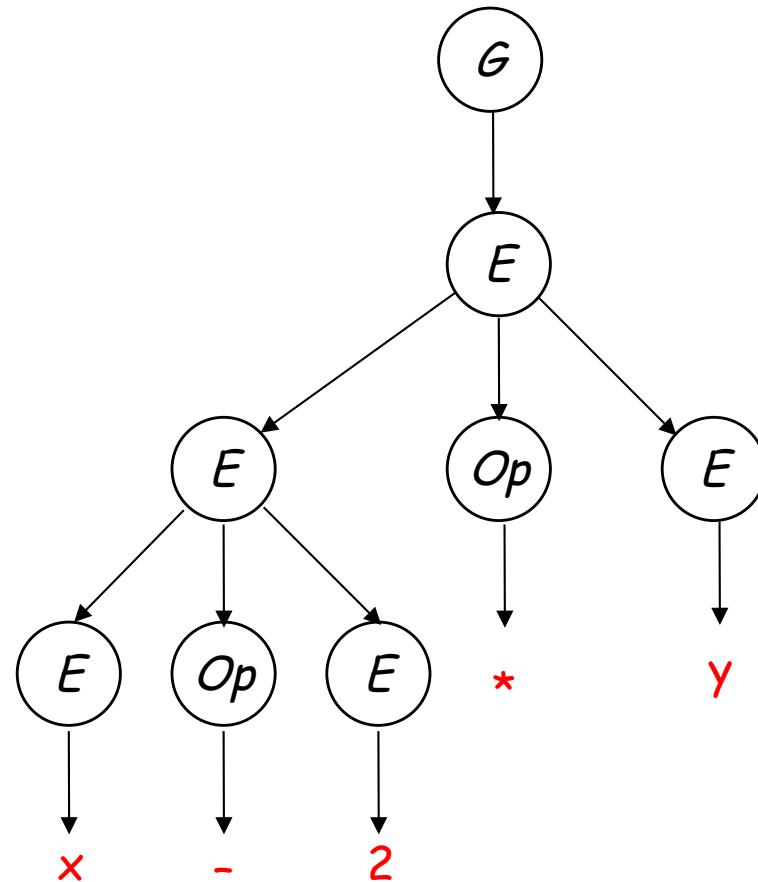


Leftmost derivation and Rightmost derivation



Leftmost derivation

This evaluates as $x - (2 * y)$



Rightmost derivation

This evaluates as $(x - 2) * y$



Derivations and Precedence

These two derivations point out a problem with the grammar:

We want same parse tree regardless of rightmost or leftmost derivation

No notion of precedence in grammar

Key: Create a non-terminal (NT) for each level of precedence



Derivations and Precedence

To add precedence

- Create a non-terminal for each *level of precedence*
- Isolate the corresponding part of the grammar
- Force the parser to recognize high precedence subexpressions first

$a + b - \underline{c * d}$

Should recognize
 $c*d$ first!

For algebraic expressions

- Multiplication and division, first *(level one)*
- Subtraction and addition, next *(level two)*



Derivations and Precedence

Adding the standard algebraic precedence produces:

level two	1	<i>Goal</i>	→	<i>Expr</i>
	2	<i>Expr</i>	→	<i>Expr + Term</i>
	3			<i>Expr - Term</i>
	4			<i>Term</i>
level one	5	<i>Term</i>	→	<i>Term * Factor</i>
	6			<i>Term / Factor</i>
	7			<i>Factor</i>
	8	<i>Factor</i>	→	<u>number</u>
	9			<u>id</u>

This grammar is slightly larger

- Takes more rewriting to reach some of the terminal symbols
- Encodes expected precedence
- Produces same parse tree under leftmost & rightmost derivations

Let's see how it parses $x - 2 * y$

Note that you can only get to Term through Expr!



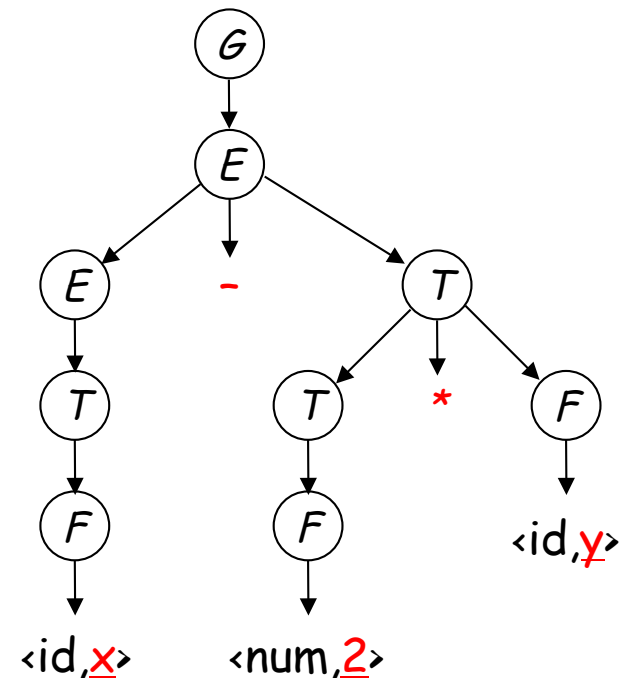
Rightmost derivation of $x-2*y$.

	1	<i>Goal</i>	→	<i>Expr</i>
level two	2	<i>Expr</i>	→	<i>Expr + Term</i>
	3			<i>Expr - Term</i>
	4			<i>Term</i>
	5	<i>Term</i>	→	<i>Term * Factor</i>
level one	6			<i>Term / Factor</i>
	7			<i>Factor</i>
	8	<i>Factor</i>	→	<u>number</u>
	9			<u>id</u>

Derivations and Precedence

Rule	Sentential Form
—	Goal
1	Expr
3	Expr - Term
5	Expr - Term * Factor
9	Expr - Term * <id, <u>y</u> >
7	Expr - Factor * <id, <u>y</u> >
8	Expr - <num, <u>z</u> > * <id, <u>y</u> >
4	Term - <num, <u>z</u> > * <id, <u>y</u> >
7	Factor - <num, <u>z</u> > * <id, <u>y</u> >
9	<id, <u>x</u> > - <num, <u>z</u> > * <id, <u>y</u> >

The rightmost derivation



Its parse tree

This produces $x - (z * y)$, along with an appropriate parse tree.

Both the leftmost and rightmost derivations give the same expression, because the grammar directly encodes the desired precedence.

Ambiguous Grammars

Our original expression grammar had other problems

- Let's look at original leftmost derivation

1	$Expr \rightarrow Expr Op Expr$
2	<u>number</u>
3	<u>id</u>
4	$Op \rightarrow +$
5	-
6	*
7	/

Rule	Sentential Form
—	$Expr$
1	$Expr Op Expr$
③	$\langle id, \underline{x} \rangle Op Expr$
5	$\langle id, \underline{x} \rangle - Expr$
1	$\langle id, \underline{x} \rangle - Expr Op Expr$
2	$\langle id, \underline{x} \rangle - \langle num, \underline{2} \rangle Op Expr$
6	$\langle id, \underline{x} \rangle - \langle num, \underline{2} \rangle * Expr$
3	$\langle id, \underline{x} \rangle - \langle num, \underline{2} \rangle * \langle id, \underline{y} \rangle$

Make note of the second rule we use!

Ambiguous Grammars

Our original expression grammar had other problems

- The grammar is *ambiguous*

1	$Expr \rightarrow Expr Op Expr$
2	<u>number</u>
3	<u>id</u>
4	$Op \rightarrow +$
5	-
6	*
7	/

Rule	Sentential Form
—	$Expr$
1	$Expr Op Expr$
①	$Expr Op Expr Op Expr$
3	$\langle id, \underline{x} \rangle Op Expr Op Expr$
5	$\langle id, \underline{x} \rangle - Expr Op Expr$
2	$\langle id, \underline{x} \rangle - \langle num, \underline{2} \rangle Op Expr$
6	$\langle id, \underline{x} \rangle - \langle num, \underline{2} \rangle * Expr$
3	$\langle id, \underline{x} \rangle - \langle num, \underline{2} \rangle * \langle id, \underline{y} \rangle$

different choice
than the first time



Two Leftmost Derivations for $x - 2 * y$

The Difference:

- Different productions chosen on the second step
- Both derivations succeed in producing $x - 2 * y$

Rule	Sentential Form
—	<i>Expr</i>
1	<i>Expr Op Expr</i>
③	$\langle \text{id}, \underline{x} \rangle \text{ Op Expr}$
5	$\langle \text{id}, \underline{x} \rangle - \text{Expr}$
1	$\langle \text{id}, \underline{x} \rangle - \text{Expr Op Expr}$
2	$\langle \text{id}, \underline{x} \rangle - \langle \text{num}, \underline{2} \rangle \text{ Op Expr}$
6	$\langle \text{id}, \underline{x} \rangle - \langle \text{num}, \underline{2} \rangle * \text{Expr}$
3	$\langle \text{id}, \underline{x} \rangle - \langle \text{num}, \underline{2} \rangle * \langle \text{id}, \underline{y} \rangle$

Original choice

Rule	Sentential Form
—	<i>Expr</i>
1	<i>Expr Op Expr</i>
①	<i>Expr Op Expr Op Expr</i>
3	$\langle \text{id}, \underline{x} \rangle \text{ Op Expr Op Expr}$
5	$\langle \text{id}, \underline{x} \rangle - \text{Expr Op Expr}$
2	$\langle \text{id}, \underline{x} \rangle - \langle \text{num}, \underline{2} \rangle \text{ Op Expr}$
6	$\langle \text{id}, \underline{x} \rangle - \langle \text{num}, \underline{2} \rangle * \text{Expr}$
3	$\langle \text{id}, \underline{x} \rangle - \langle \text{num}, \underline{2} \rangle * \langle \text{id}, \underline{y} \rangle$

New choice



Ambiguous Grammars

Definitions

- If a grammar has more than one leftmost derivation for a single *sentential form*, the grammar is *ambiguous*
- If a grammar has more than one rightmost derivation for a single sentential form, the grammar is *ambiguous*
- The leftmost and rightmost derivations for a sentential form may differ, even in an unambiguous grammar



If-then-else problem

Classic example

$Stmt \rightarrow$ if $Expr$ then $Stmt$
 | if $Expr$ then $Stmt$ else $Stmt$
 | ... other stmts ...

This ambiguity is entirely grammatical in nature

Ambiguity

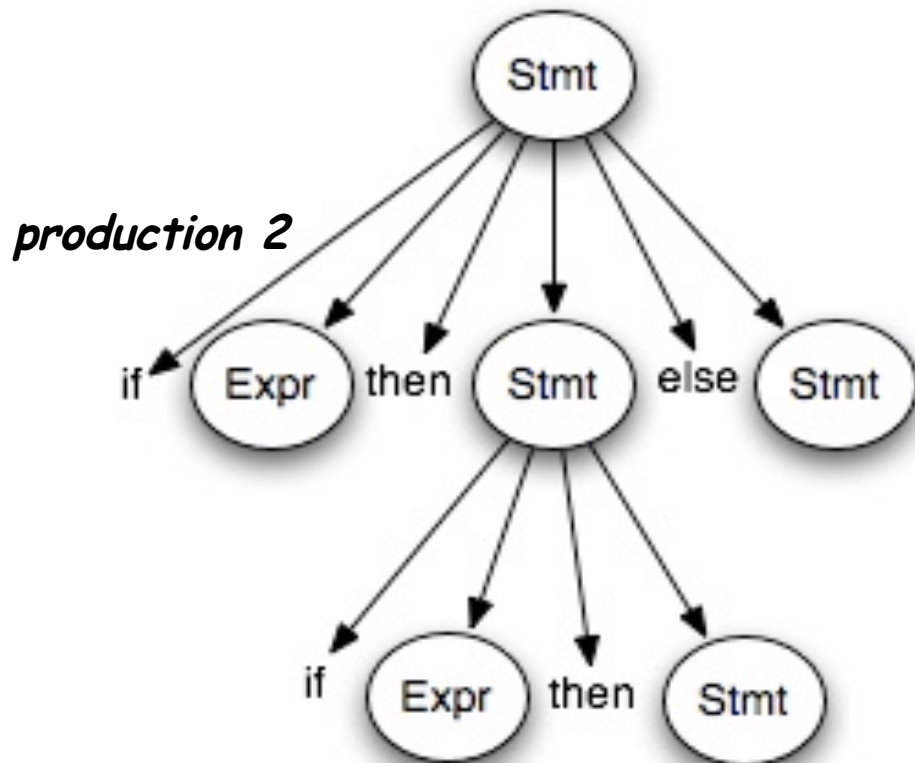
This if statement has two derivations

if $Expr_1$ then if $Expr_2$ then $Stmt_1$ else $Stmt_2$

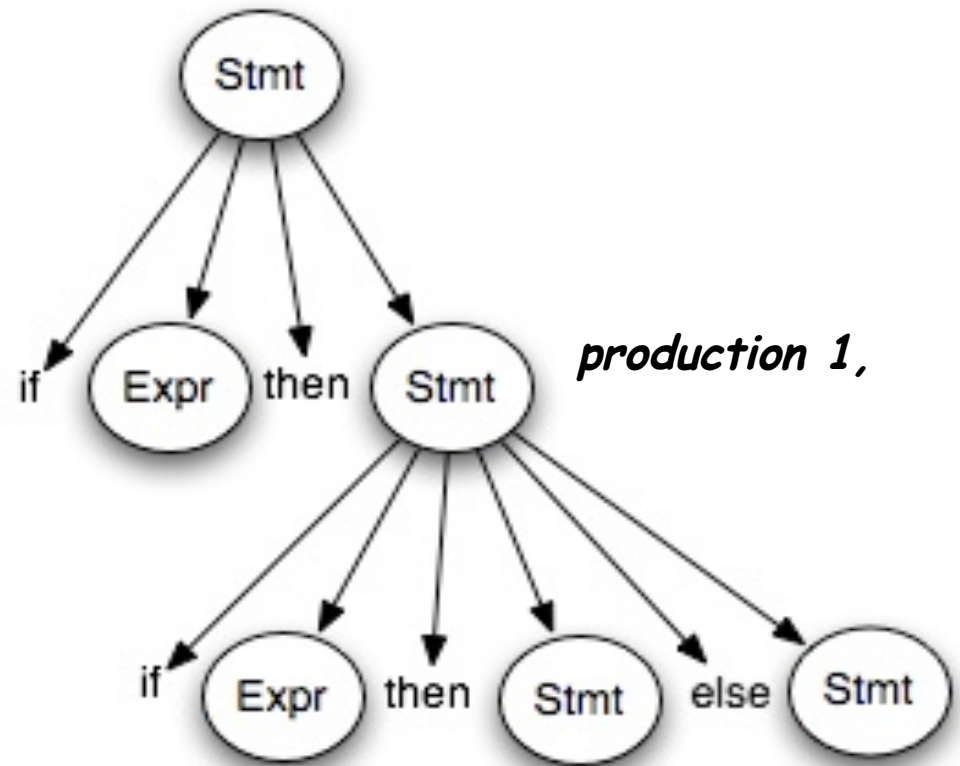
$Stmt \rightarrow$ if $Expr$ then $Stmt$ (1)

| if $Expr$ then $Stmt$ else $Stmt$ (2)

| ... other stmts ...



then production 1



production 1,

then production 2



Ambiguity

Removing the ambiguity

- Must rewrite the grammar to avoid the problem
- Match each else to innermost unmatched if (*common sense rule*)

With this grammar, the example has only one derivation

1		<i>Statement</i>	→	<u>if</u> <i>Expr</i> <u>then</u> <i>Statement</i>
2				<u>if</u> <i>Expr</i> <u>then</u> <i>WithElse</i> <u>else</u> <i>Statement</i>
3				<i>Assignment</i>
4		<i>WithElse</i>	→	<u>if</u> <i>Expr</i> <u>then</u> <i>WithElse</i> <u>else</u> <i>WithElse</i>
5				<u>Assignment</u>

Intuition: binds each else to the innermost if



Ambiguity

if $Expr_1$ then if $Expr_2$ then Assignment₁ else Assignment₂

1		<i>Statement</i>	→	<u>if</u> <i>Expr</i> <u>then</u> <i>Statement</i>
2				<u>if</u> <i>Expr</i> <u>then</u> <i>WithElse</i> <u>else</u> <i>Statement</i>
3				<i>Assignment</i>
4		<i>WithElse</i>	→	<u>if</u> <i>Expr</i> <u>then</u> <i>WithElse</i> <u>else</u> <i>WithElse</i>
5				<u>Assignment</u>

<i>Rule</i>	<i>Sentential Form</i>
	<i>Statement</i>
1	<u>if</u> <i>Expr</i> <u>then</u> <i>Statement</i>
2	<u>if</u> <i>Expr</i> <u>then</u> <u>if</u> <i>Expr</i> <u>then</u> <i>WithElse</i> <u>else</u> <i>Statement</i>
3	<u>if</u> <i>Expr</i> <u>then</u> <u>if</u> <i>Expr</i> <u>then</u> <i>WithElse</i> <u>else</u> <i>Assignment</i>
5	<u>if</u> <i>Expr</i> <u>then</u> <u>if</u> <i>Expr</i> <u>then</u> <i>Assignment</i> <u>else</u> <i>Assignment</i>

This binds the else controlling Assignment₂ to the inner if



Deeper Ambiguity

Ambiguity usually refers to confusion in the CFG

Overloading can create deeper ambiguity

$$a = f(17)$$

In many Algol-like languages, f could be either a function or a subscripted variable

Disambiguating this one requires context

- Need values of declarations
- Really an issue of *type*, not context-free syntax
- Must handle these with a different mechanism



Ambiguity - the Final Word

Ambiguity arises from two distinct sources

- Confusion in the context-free syntax (*if-then-else*)
- Confusion that requires context to resolve (*overloading*)

Resolving ambiguity

- To remove context-free ambiguity, rewrite the grammar
- To handle context-sensitive ambiguity takes cooperation
 - Knowledge of declarations, types, ...
 - Accept a superset of $L(G)$ & check it by other means[†]

[†]See Chapter 4