| Topic 2 <br> Scheme and Procedures and <br> Processes |
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|  |
| September 2008 |
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## The substitution model

- Evaluate expressions <expr1>, <expr2>, ...
- Substitute the value of <expr1> for <param1>, the value of <expr2> for <param2>, ... in a copy of the <body> expression in the definition of <fun> to make a new expression
- Evaluate that expression


## Substitution model for defined procedure application

Function definition:
(define (<fun> <param1> <param2> ...)
<body>)

Function call:
(<fun> <expr1> <expr2> ...)

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## Substitution model example

(define (square x ) (* $\mathbf{x} \mathbf{x}$ ))
Evaluation of (square 2) by substitution model:
(square 2)
(* 2 2)
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## Applicative order and normal order

- Applicative order: evaluate arguments, then apply procedure to values
- Normal order: substitute argument expressions for corresponding parameters in body of procedure definition, then evaluate body

Our substitution model of evaluation uses applicative order

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## Conditional statements

Two forms:

1) (if <test>
<then-expr>
<else-expr>) NOT optional in Scheme

## Example:

(define (absolute $x$ )
(if (< x 0)
(- x)
x) )

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## Comments on conditionals

- A test is considered to be true if it evaluates to anything except \#f
- A branch of a cond can have more than one expression:
(<test> <expr1> <expr2> . . . <exprN>)
- (<test>) returns value of <test> if it is not \#f
- The else branch must contain at least one expression


## Boolean functions

- (and <expr1> <expr2> . . . <exprN>)
<expr>s evaluated in order; return \#f if any evaluate to \#f, else return value of <exprN>
- (or <expr1> <expr2> . . . <exprN>)
<evpr>s evaluated in order; return the first value that is not \#f; return \#f if all are \#f them.
$\left.\begin{array}{l}\text { • (not <expr>) } \\ \text { returns \#t or \#f as appropriate } \\ \text { NOTE: define, if, cond, and, or are special forms } \\ \text { Fall } 2008 \\ \text { Programming Development } \\ \text { Techniques }\end{array}\right]$


## Difference between procedures and mathematical functions

Mathematical functions are defined declaratively, by stating WHAT the conditions are that their values satisfy.

Example:
$\operatorname{sqrt}(x)=$ the unique $y$ such that $y \geq 0$ and $y^{*} y=x$.

## Newton's method for square roots <br> General approach:

- If a guess is good enough, return it.
- If not good enough, compute a better guess.
- Repeat
(As a practical matter, we'll only compute an approximate value.)
To compute (squareroot $x$ ) with guess=y, new guess is $(y+x / y) / 2$
I.e., average $y$ with $x / y$

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## Square root of 10

| Guess: | Good enough? |
| :--- | :--- |
| 2 | $2 * 2=4$ |
| $(2+(10 / 2)) / 2=3.5$ | $3.5 * 3.5=12.25$ |
| $(3.5+(10 / 3.5)) / 2=$ | $3.1785 * 3.1785=$ |
| 3.1785 | 10.1029 |
| $\ldots$ | $\cdots$ |

## (code continued)

```
(define (good-enough-sqrt? guess x)
    (< (abs (- x (square guess)))
        0.000001))
(define (better-sqrt-guess guess x)
    (average guess (/ x guess)))
(define (average x y)
    (/ (+ x y) 2))
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\section*{Brain Teaser - getting recursion}
- What do the following procedures print when applied to 4 ? Note, I am not worried about the value returned, rather, about what is printed.
- First, try to predict what is printed.
- Second, try it in scheme.
- Third, if your prediction is different from what scheme produced, figure out why.
- Fourth, get help if it doesn't make sense!
(define (count1 x)
The Code
(cond ((=x0) (print x))
(else (print x)
(count1 (- x 1)))))
(define (count2 x )
(cond ((=x0) (print x\()\) )
(else (count2 (-x 1))
(print \(x)\) )))
(define (print \(x\) )
(display x )
(newline))
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\section*{Square root example}
```

Subproblems (subprocedures)
sqrt
compute-sqrt (also calls itself)
good-enough?
square
better
average
(primitives abs, /, +, - are also called)
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```


\section*{Procedural abstraction}
- A user-defined procedure is called by name, just as primitive procedures are.
- How the procedure operates is hidden.

> (square \(x\) )
> \((\exp \mathrm{x})\)

\section*{Local variables}
- The formal parameters of a procedure definition are local variables in the body.
- Other variables can become local variables by defining values for them; they become bound. variables named in the procedure head.
- Changing the name of free variables will change the meaning of the definition.
(define (area-of-circle radius)
(define pi 3.14159)
(* pi radius radius))
\begin{tabular}{|c|c|c|}
\hline \multicolumn{3}{|l|}{Procedure definitions can be nested} \\
\hline \multicolumn{3}{|l|}{\multirow[t]{2}{*}{}} \\
\hline & & \\
\hline \multicolumn{3}{|c|}{\multirow[t]{2}{*}{guess \({ }_{\text {c }}\) (cmpte-squt}} \\
\hline & & \\
\hline \multicolumn{3}{|c|}{compute-sqrt
(better guess \(x\) )} \\
\hline mil 2 nem & memempmemme & * \\
\hline
\end{tabular}

\section*{(sqrt continued)}
```

(define (good-enough? guess x)
(or (< guess 1e-100)
(< (abs (- (/ x

```
                                    (square guess))
                    1))
            \(0.000001)\) )

\section*{Block structure}
- Procedure definitions are often called blocks
- The nesting of procedure definitions is called block structure
- Variables that are free in an inner block are bound as local variables in an outer block
- Values of local variables don't have to be passed into nested definitions via parameters
- This manner of determining the values of variables is called lexical scoping

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(fewer params continued)
(define (better)
(average guess (/ x guess)))
(if (good-enough?)
guess
(compute-sqrt (better))))
(compute-sqrt 1.0))
1))
\(0.000001)\) )

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