

## A Practical Dynamic Buffer Overflow Detector

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## What's This Paper About?

- Dynamic Buffer Overflow Detectors
- CRED(C Range Error Detector)
- Results of the tests that are done by CRED.
- sOME

## Buffer Overflows

- Still being discovered in programs in use.
- In year 2002 57% of CERT advisories were related to it.
- In August 2003 50% of security advisories fell under it.

## Buffer Overflows

- Can cause billions of dollars worth damage in the computing community.
- Recent worms used this vulnerability.
- Some Examples: Slammer , CodeRed , Blaster , Welchia

## Buffer Overflow Attacks

- Attacker tries to modify the memory state to control the program in the privileged mode.
- Example:Smashing the stack.

## Smashing the Stack

- Attacker may overwrite the return address of the function on the stack.
- Attackers may modify the locations referenced by function pointer.
- Attackers may modify the global offset table

## Detecting Buffer Overflows

- A practical solution could not be found for years.
- Static Detectors vs Dynamic Detectors

## Static Buffer Overflow Detectors

- Try to verify that all memory accesses are checked for overruns.
- Disadvantages:
  - The problem is undecidable .
  - Sound tools may generate false warnings .
  - Unsound tools may miss vulnerabilities .
  - Requires manual inspection and fix of the code .

## Dynamic Buffer Overflow Detectors

- They automatically insert the necessary guards.
- It must also:
  - Protect against all buffer overflow attacks .
  - Not break working code .
  - Be efficient .

## Dynamic Buffer Overflow Detectors

- Some of them only offers protection against smash stacking.
  - Examples: StackGuard , StackShield, Propolice
- Some replaces the pointers with a structure that will be used for bounds checking.

## Dynamic Buffer Overflow Detectors(Continued)

- Jones and Kelly proposes another model by using referent objects.
- It has some weaknesses.
  - False warnings
  - Computational Overhead

## Bounds Checking Using Referent Objects

- Proposed by Jones and Kelly.
- It states that address computed from an in bound pointer must share the same referent object as the original pointer.
- It uses an object table which is stored as splay tree.

## Handling Out of Bounds Pointers

- Jones and Kelly approach pads each object by 1 byte.
- All out of bounds addresses are replaced by ILLEGAL value.
- What if we store an out of bound address, and then use it again to compute an in bound address.

## Proposed Solution for Out of Bounds Pointers(CRED)

- Create a OOB (Out of Bounds) object for every out of bounds address value.
- OOB Object includes:
  - Out of Bounds address value.
  - Referent Object that value refers to
- OOBs are stored in OOB Hash Table.

## Proposed Solution for Out of Bounds Pointers(CRED)

- Procedure
  1. If the address computed is Out of Bounds, it is replaced by OOB object and hash table is updated.
  2. When a pointer is dereferenced, if it is not pointing to an unchecked object or to an object in object table, abort and print error message.
  3. If a pointer is used in arithmetic/comparison operation, and if it is not pointing to unchecked object or to an object in the object table, value is retrieved by looking OOB hash table.

## Proposed Solution for Out of Bounds Pointers(Continued)

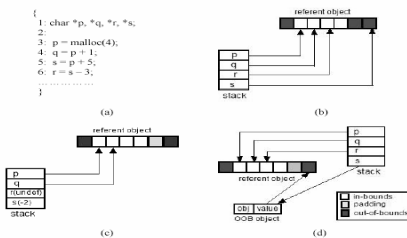
- Procedure:
  4. When an object is deleted, delete all the OOBs referring to it, update the OOB hash table.

## Example

```

1 char *p, *q, *r, *s;
2
3 p = malloc(4);
4 q = p + 1;
5 s = p + 5;
6 r = s - 3;
.....

```



## Run Time Overhead

- Performance of Jones and Kelly's technique was another disadvantage.
- Generally buffer overflow attacks uses string data.
- Also when data is copied between memory locations it's in the form of `char *`
- So bounds checking can be restricted to strings.

## Experiments with CRED

- CRED is merged with the Jones and Kelly checker for gcc 3.3.1
- It's used to test the 20 common open source programs.
- It's run by checking all buffers, not only strings

## Results of Experiment

Program	Type	# Lines	Vuln.	Tests	JK	CRED
Apache-1.3.24	web server	73.6K	no	yes	fail	pass
binutils-2.13.2.1	binary tools	596.5K	no	yes	fail	pass
bison-1.875	parser generator	25.1K	no	yes	fail	pass
ccrypt-1.4	encryption utility	4.4K	no	yes	pass	pass
coreutils-5.0	file, shell, & text utilities	69.5K	no	yes	fail	pass
encript-1.6.1	ascii to postscript converter	22.1K	no	yes	fail	pass
gawk-3.1.2	string manipulation tool	36.4K	yes	yes	fail	pass
gnupg-1.2.2	OpenPGP implementation	71.2K	no	yes	fail	pass
grep-2.5.1	pattern matching utility	20.8K	no	yes	fail	pass
gzip-1.2.4	compression utility	5.9K	yes	yes	pass	pass
hypermail-2.1.5	mail to HTML converter	27.6K	yes	yes	fail	pass
monkey-0.7.1	web server	2.5K	yes	no	pass	pass
OpenSSH-3.2.2.p1	SSH protocol implementation	43.4K	no	no	fail	pass
OpenSSL-0.9.7b	SSL & TLS toolkit	162.7K	no	yes	fail	pass
pgp4pine-1.76	mail encryption tool	3.3K	yes	no	fail	pass
polymorph-0.40	filesystem mixer	0.4K	yes	no	pass	pass
tar-1.13	archiving utility	18.2K	no	yes	pass	pass
WsMp3-0.0.10	web server	3.4K	yes	no	pass	pass
wu-ftpd-2.6.1	FTP server	18.3K	no	no	pass	pass
zlib-1.13	data compression library	8.3K	no	yes	pass	pass

## Some Comments on Results

- CRED passed all the tests.
- Jones and Kelly failed %60 of the tests.
- CRED also found some previously unknown bounds errors.

## Protection Experiments

1. 7 Known vulnerable programs are instrumented with CRED. An attacks are applied to them.
2. 20 different buffer overflows attacks are used to test the CRED.

## Results of the Performance Tests

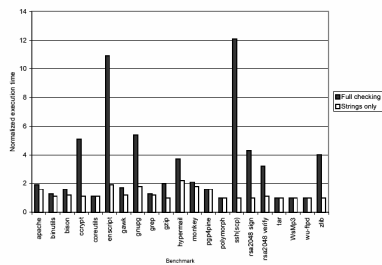
- To determine the overhead of CRED because out of bounds address tracking, we compare it with the Jones and Kelly technique

	JK(s)	CRED (s)
ccrypt	26.93	23.00
gzip-1.2.4	0.19	0.18
monkey-0.7.1	5.60	6.00
polymorph-0.4.0	0.39	0.39
tar-1.13	0.66	0.76
WsMp3-0.0.10	1.48	1.48
wu-ftpd-2.6.1	33.40	33.40
zlib-1.13	0.11	0.11

## Evaluation Criteria for 8 Programs

Program	What was evaluated
Apache-1.3.24	Response time to 15K tcp connections at the rate of 50 per second.
monkey-0.7.1	Response time to 3K tcp connections at the rate of 50 per second.
openssl-3.2.2p1	Latency of 15MB file transfer using scp via the network loop back interface.
OpenSSL-0.9.7b	Time to sign and verify 2048 bit keys using rsa.
pgp4pine-1.76	Time to decrypt 1MB file.
polymorph-0.40	Time to convert names of 100 files to unix style (lower case) names.
WsMp3-0.0.10	Latency of downloading a 15MB file.
wu-ftpd-2.6.1	Latency of 15MB file transfer via the network loop back interface.

## Performance Experiments



## Related Work

- Static Analysis Approaches
  - Wagner proposed a system where overflow detection system proposed as an integer constraint problem.
  - Larochelle and Evans presented Annotation assisted system.
  - CSSV(C String Static Verifier) detects string manipulation errors by the help of procedure summaries.

## Related Work

- Dynamic Analysis Approaches
  - StackGuard, Propolice, StackShield, Baratloo are programs that prevents stack smashing attacks.
  - Lhee and Chapin presented a buffer overflow technique using array bounds checking.
  - STOBO is a tool that can find vulnerabilities due to use of library functions.

## Related Work

- Mixed Approach
  - Ccured, Cyclone first statically detects the buffer overflows and then runs runtime tests.
  - DynamoRIO uses program shepherding to detect buffer overflows/

## Conclusions

- CRED is a practical dynamic overflow detector for C programs.
- It is built on Jones Kelly technique.
- It is more efficient and succesful with respect to other methods.
- Further improvement is possible with bounds checking and using static checking beforehand.

## Questions?